

ÉCOLE POLYTECHNIQUE FÉDÉRALE DE LAUSANNE

School of Computer and Communication Sciences

Handout 24

Solutions to Problem Set 10

Principles of Digital Communications

May 12, 2015

SOLUTION 1.

(a) In the absence of noise,

$$y(t) = \sum_{l \in \mathbb{Z}} s_l p(t - lT) \star h(t) \star q(t).$$

Because $f(t - \tau) \star g(t) = (f \star g)(t - \tau)$, we obtain

$$y(t) = \sum_{l \in \mathbb{Z}} s_l (p \star h \star q)(t - lT) = \sum_{l \in \mathbb{Z}} s_l g(t - lT).$$

(b) Following the hint, we obtain

$$\begin{aligned} g(-lT) &= \int_{-\infty}^{\infty} g_{\mathcal{F}}(f) e^{-j2\pi f lT} df \\ &= \sum_{i \in \mathbb{Z}} \int_{-\frac{1}{2T}}^{\frac{1}{2T}} g_{\mathcal{F}}\left(f - \frac{i}{T}\right) e^{-j2\pi\left(f - \frac{i}{T}\right)lT} df \\ &= \sum_{i \in \mathbb{Z}} \int_{-\frac{1}{2T}}^{\frac{1}{2T}} g_{\mathcal{F}}\left(f - \frac{i}{T}\right) e^{-j2\pi f lT} df \\ &= \int_{-\frac{1}{2T}}^{\frac{1}{2T}} \sum_{i \in \mathbb{Z}} g_{\mathcal{F}}\left(f - \frac{i}{T}\right) e^{-j2\pi f lT} df \\ &= \int_{-\frac{1}{2T}}^{\frac{1}{2T}} s(f) e^{-j2\pi f lT} df, \end{aligned}$$

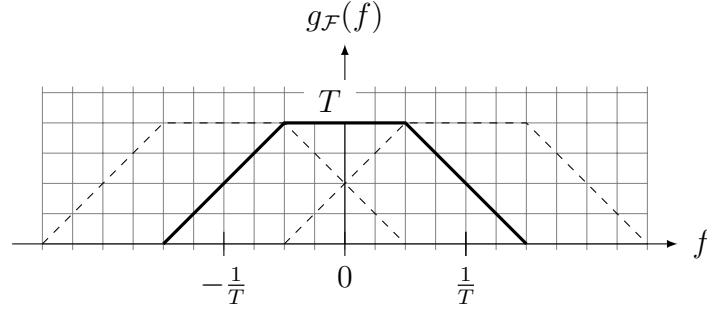
where $s(f) = \sum_{i \in \mathbb{Z}} g_{\mathcal{F}}\left(f - \frac{i}{T}\right)$.

Therefore, $g(-lT)$ are the Fourier series coefficients (multiplied by $1/T$) of the periodic function $s(f)$. If $g(lT) = \mathbb{1}\{l = 0\}$, then $s(f)$ is constant over the interval $[-\frac{1}{2T}, \frac{1}{2T}]$ and equal to T . If $s(f)$ is constant and equal to T over $[-\frac{1}{2T}, \frac{1}{2T}]$, then $g(lT) = \mathbb{1}\{l = 0\}$.

Because $s(f)$ has period $1/T$, it is equal to T for any f , not only in the interval $[-\frac{1}{2T}, \frac{1}{2T}]$.

(c) Using the hint and (b), we see that $\{\psi(t - jT)\}_{j=-\infty}^{\infty}$ is an orthonormal set if and only if $\sum_{l \in \mathbb{Z}} R_{\psi_{\mathcal{F}}}(f - \frac{l}{T}) = T$. But $R_{\psi_{\mathcal{F}}}(f) = |\psi_{\mathcal{F}}(f)|^2$, which proves Theorem 5.6.

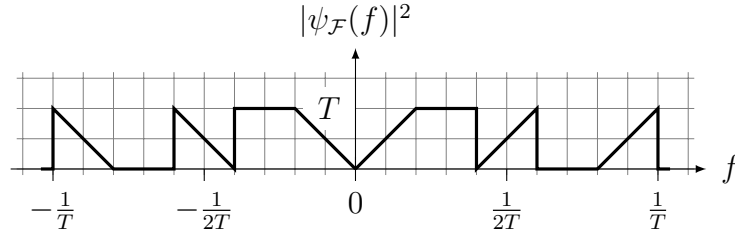
- (d) From the below figure, we see that $\sum_{i \in \mathbb{Z}} g_{\mathcal{F}}(f - \frac{i}{T}) = 2T$. Therefore, according to (b), $g(kT) = 2\mathbb{1}\{k = 0\}$ and $y(kT) = 2s_k$.



SOLUTION 2. Because $\psi(t)$ is real, its Fourier transform is conjugate symmetric ($\psi_{\mathcal{F}}(f) = \psi_{\mathcal{F}}^*(-f)$).

From the condition $\int \psi(t - kT)\psi(t - lT)dt = \mathbb{1}\{k = l\}$ for every pair k, l , it follows that $|\psi_{\mathcal{F}}(f)|^2$ satisfies Nyquist's criterion with parameter T , $\sum_{k \in \mathbb{Z}} |\psi_{\mathcal{F}}(f - k/T)|^2 = T$. On the other hand, since $\psi_{\mathcal{F}}(f) = 0$ for $|f| > \frac{1}{T}$, $|\psi_{\mathcal{F}}(f)|^2$ must have band-edge symmetry.

Putting everything together, we obtain the complete plot of $|\psi_{\mathcal{F}}(f)|^2$.



SOLUTION 3.

- (a) The self-similarity function is the inverse Fourier transform of $|\psi_{\mathcal{F}}(f)|^2$, hence

$$|\psi_{\mathcal{F}}(f)|^2 = T\mathbb{1}\{f \in [-1/2T, 1/2T]\} \iff R_{\psi}(\tau) = \text{sinc}(\tau/T) = \frac{\sin(\pi\tau/T)}{\pi\tau/T}$$

To determine the time-domain pulse $\psi(t)$, we first observe that the frequency domain pulse can be

$$\psi_{\mathcal{F}}(f) = \sqrt{T}\mathbb{1}\{f \in [-1/2T, 1/2T]\}e^{-j2\pi ft_0}, \quad \text{for } \forall t_0 \in \mathbb{R}.$$

Consequently,

$$\psi(t) = \frac{1}{\sqrt{T}} \text{sinc}\left(\frac{t - t_0}{T}\right) \quad \text{for } \forall t_0 \in \mathbb{R}.$$

(b) We have

$$\begin{aligned}
Y_j &= \int_{-\infty}^{+\infty} R(t)\psi^*(t-jT) dt \\
&= \int_{-\infty}^{+\infty} W(t)\psi^*(t-jT) dt + \underbrace{\int_{-\infty}^{+\infty} N(t)\psi^*(t-jT) dt}_{\triangleq Z_j} \\
&= \int_{-\infty}^{+\infty} \sum_i S_i \psi(t-iT)\psi^*(t-jT) dt + Z_j \\
&= \sum_i S_i \int_{-\infty}^{+\infty} \psi(t-iT)\psi^*(t-jT) dt + Z_j \\
&= \sum_i S_i \int_{-\infty}^{+\infty} \psi(t-iT)\psi^*(t-jT) dt + Z_j \\
&= \sum_i S_i R_\psi((j-i)T) + Z_j. \tag{1}
\end{aligned}$$

Since $N(t)$ is a white Gaussian noise we know Z_j are zero-mean jointly Gaussian random variables with covariance

$$\text{cov}(Z_i, Z_j) = \frac{N_0}{2} \int \psi^*(t-iT)\psi(t-jT) dt = \frac{N_0}{2} R_\psi((i-j)T). \tag{2}$$

From Problem 1 we know that $\psi(t)$ is orthonormal to its time translates by T , which is equivalent to

$$R_\psi(kT) = \mathbb{1}\{k=0\}.$$

Using the above in (1) and (2) shows that

$$Y_j = S_j + Z_j,$$

where Z_j are i.i.d. $\mathcal{N}(0, \frac{N_0}{2})$ noise samples.

(c) Following the same steps as in (b) we have

$$\begin{aligned}
Y_j &= \int_{-\infty}^{+\infty} R(t)\psi^*(t - jT - \Delta) dt \\
&= \int_{-\infty}^{+\infty} W(t)\psi^*(t - jT - \Delta) + \underbrace{\int_{-\infty}^{+\infty} N(t)\psi^*(t - jT - \Delta) dt}_{\triangleq Z_j} dt \\
&= \int_{-\infty}^{+\infty} \sum_i S_i \psi(t - iT)\psi^*(t - jT - \Delta) dt + Z_j \\
&= \sum_i S_i \int_{-\infty}^{+\infty} \psi(t - iT)\psi^*(t - jT - \Delta) dt + Z_j \\
&= \sum_i S_i \int_{-\infty}^{+\infty} \psi(t - iT)\psi^*(t - jT - \Delta) dt + Z_j \\
&= \sum_i S_i R_\psi((j - i)T + \Delta) + Z_j. \tag{3}
\end{aligned}$$

Once more since $N(t)$ is a white Gaussian noise we know Z_j are zero-mean jointly Gaussian random variables with covariance

$$\text{cov}(Z_i, Z_j) = \frac{N_0}{2} \int \psi^*(t - iT - \Delta)\psi(t - jT - \Delta) = \frac{N_0}{2} R_\psi((i - j)T),$$

which shows Z_j are still i.i.d. $\mathcal{N}(0, \frac{N_0}{2})$ noise samples. However, the samples are now of the form

$$Y_j = l_0 S_j + \sum_{i \neq j} S_i l_{j-i} + Z_j$$

where

$$l_k = R_\psi(kT + \Delta) = \text{sinc}\left(k + \frac{\Delta}{T}\right). \tag{4}$$

(d) Suppose $\Delta = \frac{T}{2}$. Then,

$$l_k = \text{sinc}\left(k + \frac{1}{2}\right) = \frac{\sin(k\pi + \pi/2)}{k\pi + \pi/2} = \frac{(-1)^k}{k\pi + \pi/2}.$$

Now if s_1, \dots, s_n is a sequence satisfying $s_i = -s_{i+1}$, for $i = 1, 2, \dots, n - 1$ and arbitrary $s_n \in \{\pm 1\}$, we have

$$\begin{aligned}
Y_n &= \frac{2}{\pi} s_n + \sum_{k \neq 0} s_{n-k} l_k + Z_n \\
&= \frac{2}{\pi} s_n + \sum_{k=1}^n s_{n-k} l_k + Z_n \\
&= \frac{2}{\pi} s_n + \sum_{k=1}^n (-1)^k \frac{(-1)^k}{k\pi + \pi/2} s_n + Z_n
\end{aligned}$$

The ISI term is hence

$$\begin{aligned} \sum_{k=1}^n (-1)^k \frac{(-1)^k}{k\pi + \pi/2} s_n &= \sum_{k=1}^n \frac{1}{k\pi + \pi/2} s_n \\ &= s_n \left(\sum_{k=0}^n \frac{1}{k\pi + \pi/2} - \frac{\pi}{2} \right) \end{aligned}$$

We see that the term inside the parenthesis diverges as n grows large.

SOLUTION 4.

(a) When $i = j$, $\mathbb{E}[X_i X_j]$ is

$$\mathbb{E}[X_i^2] = \mathbb{E}[1] = 1.$$

Remember that the B_i are iid Bernoulli($\frac{1}{2}$) random variables. Hence, we find immediately

$$\begin{aligned} K_X[1] &= \mathbb{E}[X_{2n} X_{2n+1}] = \mathbb{E}[B_n B_{n-2} B_n B_{n-1} B_{n-2}] \\ &= \mathbb{E}[B_n^2 B_{n-1} B_{n-2}^2] \\ &= \mathbb{E}[B_{n-1}] = 0, \end{aligned}$$

and also

$$\begin{aligned} K_X[2] &= \mathbb{E}[X_{2n} X_{2n+2}] = \mathbb{E}[B_n B_{n-2} B_{n+1} B_{n-1}] \\ &= \mathbb{E}[B_n] \mathbb{E}[B_{n-2}] \mathbb{E}[B_{n+1}] \mathbb{E}[B_{n-1}] = 0. \end{aligned}$$

By continuing this argument we find

$$K_X[i] = \mathbb{1}\{i = 0\}.$$

Hence,

$$S_X(f) = \frac{\mathcal{E}_s}{T_s} |\psi_{\mathcal{F}}(f)|^2.$$

This means that by choosing $\psi(t)$ appropriately, we can control the bandwidth consumption of our communications scheme.

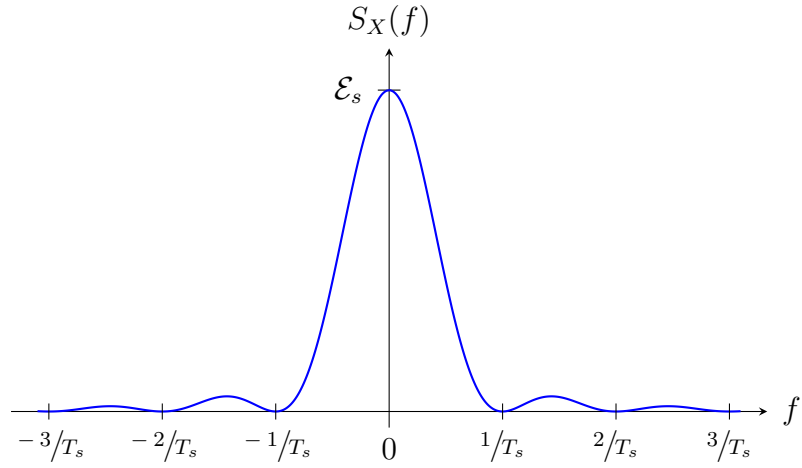
(b) We know that

$$|\psi_{\mathcal{F}}(f)|^2 = T_s \operatorname{sinc}^2(T_s f).$$

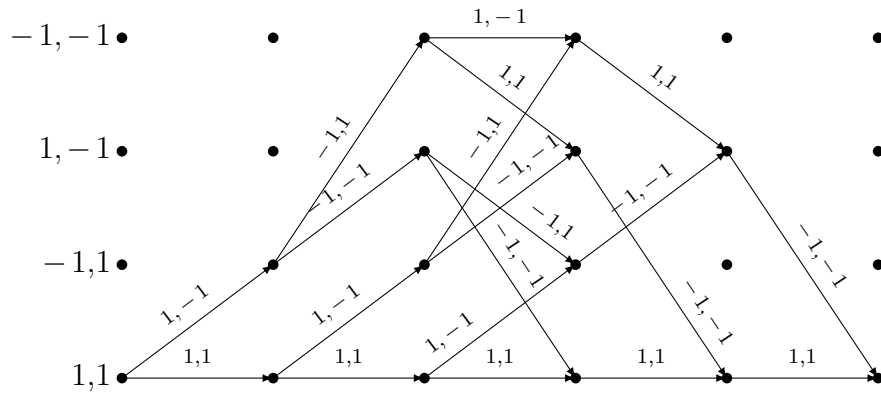
It follows that

$$S_X(f) = \mathcal{E}_s \operatorname{sinc}^2(T_s f).$$

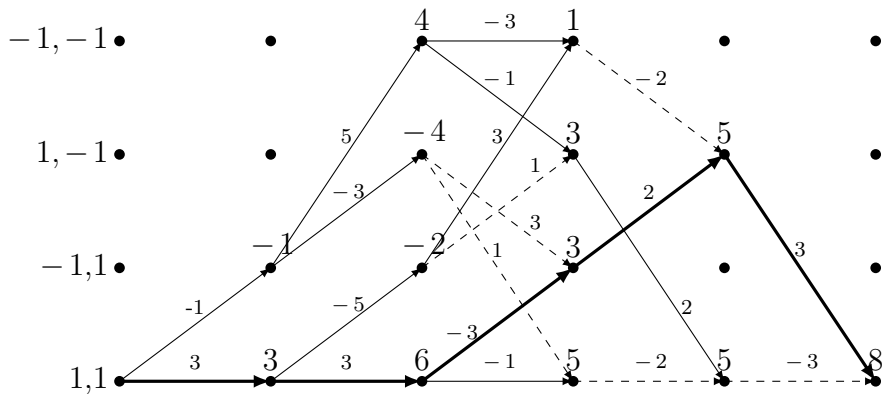
A plot of $S_X(f)$ is shown here:



SOLUTION 5. The trellis representing the encoder is shown below:



We display the diagram labeled with edge-metric according to the received sequence and state-metric of the survivor path. We also indicate the survivor paths and the decoding path.



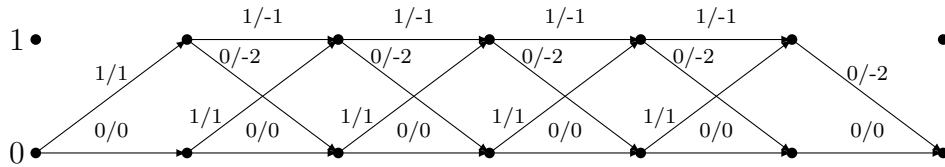
From the figure we can read the decoded sequence 1, 1, -1, 1, 1.

SOLUTION 6.

(a)

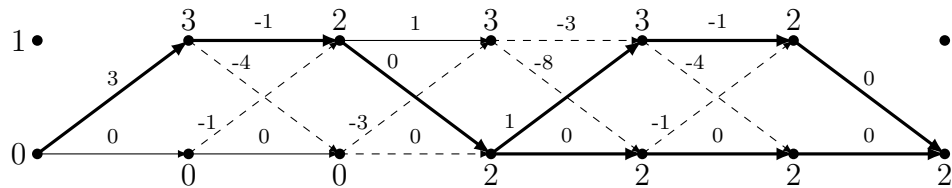
$$X_i = B_i - 2B_{i-1}$$

From this, we can draw the following trellis:



(b) We have $Y = X + Z$, where $Z = (Z_1, \dots, Z_6)$ is a sequence of iid components with $Z_i \sim \mathcal{N}(0, \sigma^2)$. Our maximum likelihood decoder is a minimum distance decoder. We have to minimize $\|y - x\|^2$ or equivalently, maximize $2\langle y, x \rangle - \|x\|^2$. We thus have $f(x, y) = \sum_{i=1}^6 2y_i x_i - x_i^2$ whose maximization with respect to x leads to a maximum likelihood decision on X .

(c) We label our trellis with the edge metric $2y_i x_i - x_i^2$ and then trace back the decoding path.



We see that the two sequences 1, 1, 0, 0, 0 and 1, 1, 0, 1, 1 are equally likely, so the decoder would choose either of the two.